

# Elliptic Curve Arithmetic

## Why ECC?

- ✓ Good security even with far smaller key than RSA
- ✓ Reduce the processing overhead

## Basics for ECC:

### Abelian Groups:

- ✓  $\{G, \bullet\}$  is a set of elements with a binary operations
- ✓ It should satisfy the following properties
  - Closure : if  $a$  and  $b$  belongs to  $G$ , than  $a \bullet b$  also in  $G$ .
  - Associative :  $a \bullet (b \bullet c) = (a \bullet b) \bullet c$  for all  $a, b, c$  in  $G$ .
  - Identity elements:  $a \bullet e = e \bullet a = a$  for all  $a$  in  $G$ .
  - Inverse element :  $a \bullet a' = a' \bullet a = e$ .
  - Commutative:  $a \bullet b = b \bullet a$  for all  $a, b$  in  $G$ .
- ✓ In Diffie-Hellman, keys are generated by exponentiation defined by repeated multiplication
  - $a^k \text{ mod } q = (a \times a \times a \times \dots \times a) \text{ mod } q$
- ✓ In ECC, operations are in addition. Multiplication is defined by repeated addition.
  - $a \times k = (a + a + a + \dots + a)$

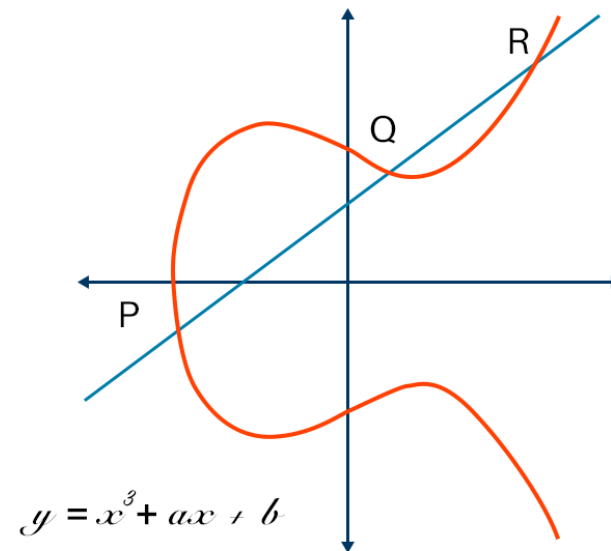
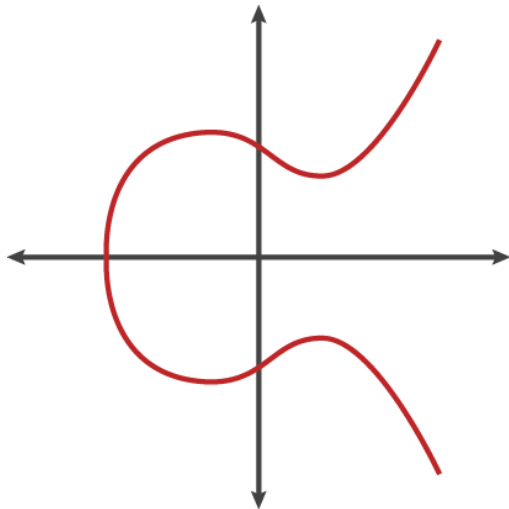
## ECC definition

- ❖ Elliptic curve is represented as  $E_p(a,b)$ .  $P$  is a prime number and  $a,b$  are restricted to mod  $p$
- ❖ The curve is represented by  $y^2 = x^3 + ax + b$ .
- ❖ Elliptic curves are not ellipse but the equation of ecc is described by calculation of circumference of ellipse (i.e. cubic equation with highest degree of 3)
- ❖ For determining the security of various elliptic curve ciphers, the number of points in a finite abelian group defined over an elliptic curve.
- ❖ In case of the finite group  $E_p(a,b)$ , the number of points  $N$  is bounded by:
$$p + 1 - \sqrt{p} \leq N \leq p + 1 + 2\sqrt{p}$$
- ❖ No. of points in  $E_p(a,b)$  is approximately equal to the number of elements in  $Z_p$ , namely  $p$  elements.

# How the elliptic curve is Symmetric?

$$y = \sqrt{x^3 + ax + b}$$

This gives a value of  $x$  as  $\pm x$ . So, each curve is symmetric curve about  $y=0$ .



Affine points: The points present in Elliptic curve

O points: There is a point called 0 point in which  $P + (-P)$  becomes infinity.

ECC can be defined as EC over  $Z_p$  (prime curve) and EC over  $GF(2^m)$  (binary curve).

ECC can be used for key exchange and Encryption.

## Elliptic curves over $Z_p$ :

- ✓ The curve of this type is prime curve
- ✓ The variables and coefficients are restricted to elements of a finite field
- ✓ The values are restricted from 0 through  $p-1$ , If the values exceeds the range perform modulo  $p$ .
- ✓ The curve is represented by  $y^2 \bmod p = (x^3+ax+b) \bmod p$
- ✓ The curve is to be focused in only one of the quadrant from  $(0,0)$  through  $(p-1,p-1)$  containing non negative integers
- ✓ The number of points  $N$  is bounded by
$$p + 1 - \sqrt{p} \leq N \leq p + 1 + 2\sqrt{p}$$

# Addition

✓ Adding 2 points  $P(x_p, y_p)$  and  $Q(x_q, y_q)$  gives  $R(x_r, y_r)$ .

✓ Steps:

✓ Find the slope  $\lambda$ :

✓  $\lambda = (y_q - y_p) / (x_q - x_p)$  if  $P \neq Q$

✓  $\lambda = (3x_p^2 + a) / 2y_p$  if  $P = Q$  where  $a$  is obtained from  $E_p(a, b)$

✓ Find the sum:  $R$  (i.e.  $(x_r, y_r)$ ) =  $P + Q$

✓  $x_r = \lambda^2 - x_p - x_q$

✓  $y_r = \lambda(x_p - x_r) - y_p$

Negating a point:

- ✓ if  $Q = (x_q, y_q)$
- ✓ then  $-Q = -(x_q, y_q) = (x_q, -y_q)$

Subtraction:

- ✓  $P - Q = (x_p, y_p) - (x_q, y_q) = (x_p, y_p) + (x_q, -y_q \text{ mod } p)$ .  
Now perform addition.

Multiplication:

- ✓ Only Scalar multiplication is possible.  
Multiplication between two points are not possible.  
Repeated addition is performed.
- ✓  $2P = P+P$ ,  $3P = P+P+P$  and so on. Note for slope ( $\lambda$ ) calculation use the formula  $P=Q$ .
- ✓ Division: Only Scalar division is possible.  $[1/a(x_p, y_p)] = a^{-1} (x_p, y_p)$ . Multiplication steps can be followed.